# 2. Exercise "Datenstrukturen und Effiziente Algorithmen", WS 18/19

## **Exercise 1:** (2.5+2.5=5 Credits)

Let B a decision problem in the class  $\mathcal{P}$  and A a decision problem with  $A \leq_p B$ .

- (a) Explain in a few words if  $A \notin \mathcal{NP}$  or  $A \in \mathcal{NP}$ .
- (b) What can you conclude about the "subset relation" of the classes  $\mathcal{P}$  and  $\mathcal{NP}$  if Problem A is  $\mathcal{NP}$ -hard. Explain your results.

## Exercise 2: "Subgraph Isomorphism" (10 Credits)

Show that the following decision problem is NP-complete:

Given two undirected graphs H and G.

Is there an isomorphism  $\varphi$  from H to some subgraph  $G' \subseteq G$ ?

HINT: You may start the reduction from CLIQUE.

## Exercise 3: "BoolMatrixDecomposition" (15 Credits)

The following problem, called SETBASISPROBLEM, is NP-complete:

Given a set C of subsets of a finite set U and a positive integer k.

Is there a set  $\mathcal{B}$  of subsets of U with  $|\mathcal{B}| \leq k$  such that for every  $C \in \mathcal{C}$  there is a subset  $\mathcal{B}_C \subseteq \mathcal{B}$  with  $\bigcup_{B \in \mathcal{B}_C} B = C$ ?

Now, let A and B be binary matrices of dimensions  $n \times k$  and  $k \times m$ , respectively. Let  $C = A \circ B$  denote the  $(n \times m \text{ matrix})$  Boolean product of A and B, i.e., the usual matrix product with the exception that 1 + 1 = 1. By way of example

$$\begin{pmatrix} 1 & 1 & 0 \\ 1 & 1 & 1 \\ 0 & 1 & 1 \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ 1 & 1 \\ 0 & 1 \end{pmatrix} \circ \begin{pmatrix} 1 & 1 & 0 \\ 0 & 1 & 1 \end{pmatrix}.$$

Use reduction from SETBASISPROBLEM to show that the following problem BOOLMA-TRIXDECOMPOSITION is NP-complete.

Given an  $n \times m$  matrix C and a positive integer  $k < \min\{m, n\}$ . Is there a  $n \times k$  matrix A and a  $k \times m$  matrix B such that  $C = A \circ B$ ?

### Deadline: Wednesday - October 31, 2018 - 12.15pm