# Bioinformatics 

(Intro Graphs)

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## Graphs and some (Basic-)Properties

- simple (un-)directed graphs $G=(V, E)$
- degree $\operatorname{deg}(v)$

Theorem
For any graph $G=(V, E)$ holds: $\sum_{v \in V} \operatorname{deg}(v)=2|E|$.
Theorem
Every graph contains an even number of vertices with odd degree.

## Graphs and some (Basic-)Properties

- $\simeq$ isomorphism
- (induced) subgraph
- complete graphs $K_{n}$


## Graphs and some (Basic-)Properties

- walk, trail, path, cycle
- connectedness


## Graphs and some (Basic-)Properties

- forests and trees


## Theorem

Let $G=(V, E)$ be an undirected graph. $G$ is a tree if and only if for any two vertices $u, v \in V$ there is exactly on path connecting them.

## Theorem

Let $G=(V, E)$ be an undirected connected graph.
$G$ is a tree if and only if $|E|=|V|-1$.
Corollary
Let $G=(V, E)$ be a connected undirected graph.
Then $|E| \geq|V|-1$ and $G$ has a spanning tree.

## Graphs and some (Basic-)Properties

- distances


## Lemma (Subpaths)

Any connected subgraph of a shortest path is a shortest path.
Lemma (Triangle Inequality)
Let $G=(V, E)$ be an undirected graph. For all $u, v, z \in V$ holds that

$$
d(u, v) \leq d(u, z)+d(z, v)
$$

## Breadth-first search - BFS (G, s)

1: INPUT: Adjacencylist of a graph $G=(V, E)$, source vertex $s$
2: for all $v \in V$ do
3: $\quad \delta[v]=\infty, \operatorname{pred}[v]=\mathrm{NIL}$
end for
5: $\delta[s]=0$, enqueue $(Q, s)$
while $Q \neq \emptyset$ do
7: u=dequeue(Q)
8: $\quad$ for all $v \in \operatorname{adj}(u)$ do
9: $\quad$ if $\delta[v]=\infty$ then
10: $\quad \delta[v]=\delta[u]+1$
11: $\quad \operatorname{pred}[v]=u$
12: enqueue( $\mathrm{Q}, \mathrm{v}$ )
13: end if
14: end for
15: end while
16: OUTPUT: Arrays $\delta$, pred;

## Breadth-first search - BFS (G, s)

## Theorem

$\operatorname{BFS}(G, s)$ runs in time $O(|V|+|E|)$.
Theorem
After termination of BFS $(G, s)$ we have $\delta(v)=d_{G}(s, v)$ for all $v \in V$

## Graphs and Colorings

- (proper) coloring
- k-coloring and chromatic number $\chi(G)$

Theorem
For all graphs $G$ holds: $\chi(G) \leq 1+\Delta(G)$.

## Graphs and Colorings

- bipartite graphs

Lemma
A graph $G$ is bipartite if and only if $G$ is 2-colorable.
Theorem
A graph $G$ is bipartite if and only if $G$ does not contain odd cycles.

